

Final Review

CSCE 742 - Lecture 24 - 12/06/2018

The Final...

- 150 minutes: December 11, 9:00 - 11:30
- In this room!
- Closed book, no notes!
- Graded on quality of answers, not how much you wrote.
- Very similar in format to the practice final and midterm.
- Not cumulative!
- Study your homework.

The Final

- Topics:
 - Concurrency Viewpoint
 - Development Viewpoint
 - Deployment Viewpoint
 - Operational Viewpoint
 - Performance and Scalability Perspective
 - Security Perspective
 - Availability Perspective
 - Real-time systems
 - Architectural Description Languages
 - Services and Service-Oriented Architecture
 - REST/SOAP
 - ML Systems

General Questions

- Today: Go over practice final questions.
- First - any general questions on course content or homework?

Question 1

Speculate as to why none of the architecture description languages other than UML have achieved widespread use.

Question 1

Speculate as to why none of the architecture description languages other than UML have achieved widespread use.

Many possible answers here. Some starting points:

- Difficult to create critical mass to attract tool builders.
- Most notations only capture a small portion of architectural concerns.
- Lack of communication between “academic” software engineering and commercial SE.
- Too domain specific; have not evolved to describe new kinds of architectures (e.g. web services).
- Too hard to use / not understandable to developers.

Question 2

Create an attack tree describing how an attacker might attempt to steal money from an Automated Teller Machine (ATM).

Attack Trees

- Structured notation for categorizing threats and their probability.
 - Represented visually as a tree as a nested list.
 - Root of the tree shows the goal of the attack.
 - Branches classify the different types of attacks that could be attempted.
 - Create a tree for each goal an attacker may have. Can be used to analyze security policies.

Goal: Obtain customer credit card details.

1. Extract details from database.
 - 1.1 Access database directly.
 1. Crack/guess database passwords.
 2. Crack/guess OS passwords that bypass database security.
 - 1.2 Access via a member of the administration staff.
 1. Bribe a database administrator (DBA).
 2. Conduct social engineering by e-mail to trick the DBA into revealing details

Attack Tree Example

2. Extract details from Web interface.
 - 2.1. Set up a dummy Web site and e-mail users the URL to trick them into entering credit card details.
 - 2.2. Crack/guess passwords for user accounts and extract details from the GUI.
 - 2.3. Send users a program by e-mail to record keystrokes.
 - 2.4. Attack the domain name server to hijack domain name and attack 2.1.
 - 2.5. Attack the server software directly to try to find loopholes in its security.
3. Find details outside the system.
 - 3.1. Conduct social engineering by phone/e-mail to get customer services staff to reveal card details.
 - 3.2. Direct a social-engineering attack on users by using public details from the site to make contact.

Question 2

Create an attack tree describing how an attacker might attempt to steal money from an Automated Teller Machine (ATM).

Goal: Steal money from ATM

1. Physical attack
 - a. Break ATM casing and steal money
 - i. [AND]
 1. Steal entire ATM for later dismantling
 2. Procure vehicle capable of transporting ATM
 - b. Card-data stealing attack
 - i. Card-based attack using new cards
 1. [AND]
 - a. Buy/steal card producing device
 - b. Buy/steal card stock for new ATM cards
 - ii. Data capture for ATM spoofing
 1. Capture ATM track 1&2 data and valid PIN using a skimming device
 2. Buy data and PINs from black market
2. Capture/guess ATM data from online banking site
3. Capture/guess ATM data through ATM software vulnerability
4. Get valid card in someone else's name (id theft)
5. ...

Question 3

What difficulties do cyclic component dependencies lead to in an architecture?

What can be done to break cyclic dependencies?

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What difficulties do cyclic component dependencies lead to in an architecture?

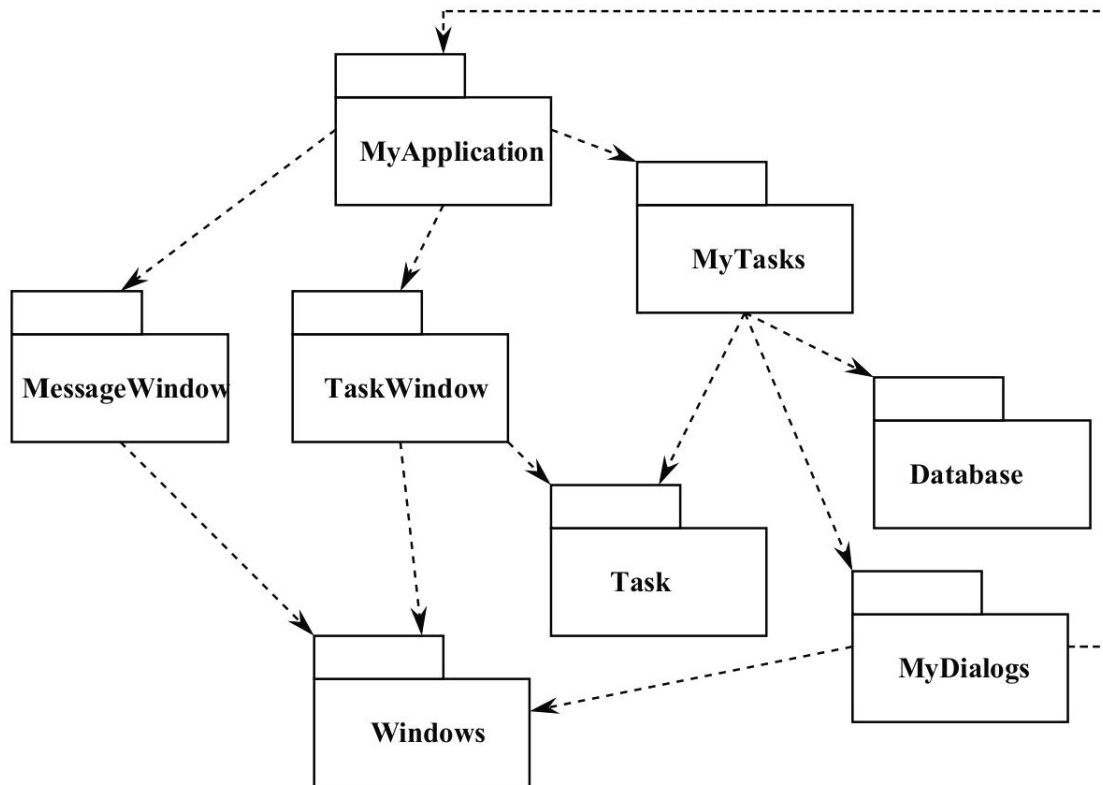
- Components that contain circular dependencies are much more difficult to test and maintain.
 - They are tightly coupled to the interface of the others.
 - For testing, a usual strategy is to use “leveling” to test.
 - Test components that have no dependencies first, then once you have confidence in them, test components that depend only on those components.
 - If you have circular dependencies, it is much more difficult to test in this way. We have to examine the behavior of all of these components simultaneously.

Question 3

What difficulties do cyclic component dependencies lead to in an architecture?

- For maintenance, it is often difficult to modify one of the cyclically connected components without changing all of the components in the cycle.
 - This can be problematic if the components containing the cyclic references reside in multiple packages.
 - For versioning, if components mutually depend then they must be installed and updated simultaneously.
- The same thing is true for compilation. A change in one triggers compilation of all others.

Cyclic Example (Bad)



Release:

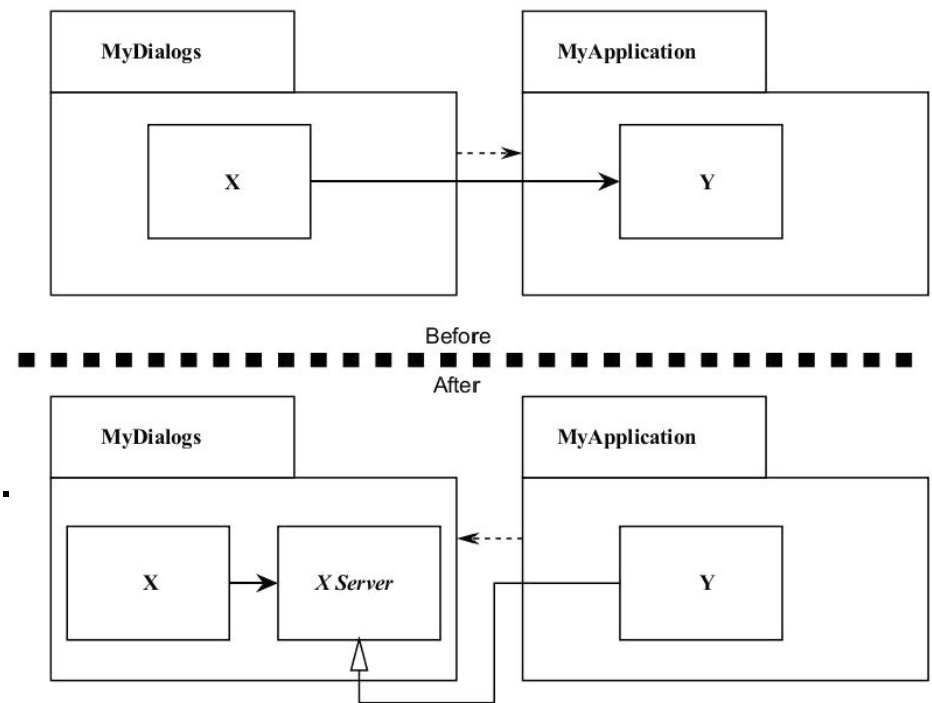
- Must be simultaneous with MyApplication.
- But this means MyTasks must also be coordinated (it is a dependency of MyApplication and depends on MyDialogs).
- This means that it must also be coordinated with Task and Database (dependencies of MyTasks).

Testing:

- MyDialogs requires MyApplication, so...
- **MyDialogs is dependent on all packages(!) for testing!**

Fixing Circular Dependencies

- Apply the Dependency Inversion Principle.
 - Create an abstract class with the interface MyDialogs needs.
 - Put the class into MyDialogs
 - Inherit into MyApplication.
 - Reverses the dependency, breaking the cycle.



Package Refactoring

- Can also create third package with class(es) that both MyApplication and MyDialog depend on.
- Package contents and dependency hierarchy must be actively managed and refactored.
 - Dependencies will change as system expands.
 - Circular dependencies must be pruned out.
 - Coordinating this movement is important job for architect.

Question 4

RPCs and messaging schemes are inter-process communication mechanisms. For one process with multiple threads:

- 1. Are there analogous concepts to RPCs and messaging between threads?**
- 2. Describe an additional means of communication that is available between threads.**
- 3. Do the same benefits/drawbacks between RPCs and messaging schemes exist when considering inter-thread communication as interprocess communication?**

Question 4

Are there analogous concepts to RPCs and messaging between threads?

- **Messaging:**
 - Constructed from semaphores (limits number of consumers for a resource).
 - Threads communicate through event queues (send event to queue).
 - Delivery is guaranteed, as we act within one process.
- **RPC:**
 - No immediately analogous concept.
 - Local procedure calls are similar, but not a means of cross-thread communication.

Question 4

Describe an additional means of communication that is available between threads.

- **Shared memory**
 - Threads can read and write from the same memory space. Form of message passing, indirectly causes changes to occur in other threads.
- **Monitors**
 - In Java, objects have a monitor that ensures that only one thread can execute a critical section of code at a time.
 - **wait()** tells calling thread to give up monitor and go to sleep until some other thread enters the same monitor and call notify.
 - **notify()** wakes up a thread that called wait() on same object.
 - **notifyAll()** wakes up all the threads that called wait() on same object.

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- **Monitors**

- wait() tells calling thread to give up monitor and go to sleep until some other thread enters the same monitor and call notify.
- notify() wakes up a thread that called wait() on same object.
- notifyAll() wakes up all the threads that called wait() on same object.
 - Threads enter to acquire lock.
 - Lock is acquired by one thread.
 - Now, thread goes to waiting state if you call wait() method on the object. Otherwise it releases the lock and exits.
 - If you call notify() or notifyAll() method, thread moves to the notified state (runnable state).
 - Now thread is available to acquire lock.
 - After completion of the task, thread releases the lock and exits the monitor state of the object.

Question 4

Describe an additional means of communication that is available between threads.

- Semaphores
 - Counter that acts as a permit for a shared resource.
 - When a thread gets a permit, the counter goes down.
 - Otherwise, the thread will be blocked until it can get a permit.
 - When finished, the thread releases the semaphore
 - (counter increments).
 - Semaphores can be used to control synchronization between threads.

Question 4

Do the same benefits/drawbacks between RPCs and messaging schemes exist when considering inter-thread communication as interprocess communication?

- RPCs may not be useful between threads. They share an address space. Why not just use the active thread?
 - No performance benefit since RPCs are synchronous.
- Messaging is still useful.
 - Allows decoupling of processes from UI, etc.
 - Allows a clean interface between threads, preventing deadlock/race conditions.
 - No guarantee of persistence (like inter-process communication), but no process lifecycle issues.

Question 4

Do the same benefits/drawbacks between RPCs and messaging schemes exist when considering inter-thread communication as interprocess communication?

- **Performance:**
 - Threads communicate faster than processes, so lower cost of splitting tasks.
- **Isolation:**
 - Threads are not isolated (like processes), so some benefits of IPC do not impact inter-thread communication.
- **IPCs allow processes to be spread across multiple machines. ITC does not.**
 - No scalability benefit.

Question 5

What are the benefits and drawbacks of using XML vs. binary protocols for messaging between processes?

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What are the benefits and drawbacks of using XML vs. binary protocols for messaging between processes?

- Custom binary protocol may require less network bandwidth than XML.
- If processes can agree on format, binary protocol may require less translation and parsing at network boundaries.
 - For applications that are network intensive, this may lead to better performance and network utilization.

Question 5

What are the benefits and drawbacks of using XML vs. binary protocols for messaging between processes?

- XML will be easier to maintain. Can use standard tools to process message traffic and diagnose errors.
- XML supports transparent re-hosting on different platforms.
- XML may lead to better performance given static data. Can perform pre-built caching.
- Can more easily add additional data to XML messages without disrupting existing clients.
- Several tools for manipulating and routing XML messages, making it easier to allow new applications into a system.

Question 6

Big Bang, Parallel Run, and Staged Migration are techniques for upgrading existing software installations.

- 1. Briefly explain each technique.**
- 2. Describe advantages and disadvantages for each technique in terms of data migration, complexity, and rollback in case of failure.**
- 3. Describe installation scenarios where you might use each technique.**

Question 6

Briefly explain each technique.

- **Big Bang:** Pick a day, turn the old system off and the new system on.
- **Parallel Run:** Choose a period in which the old and new systems run in parallel.
- **Staged Migration:** Swap out pieces of the old system with pieces of the new system over an installation period or migrate portions of the organization over a period.

Question 6

Describe advantages and disadvantages for each technique in terms of data migration, complexity, and rollback in case of failure.

- Big Bang:
 - Data migrates in one shot.
 - Simple in the sense that only one system runs at a time. Hard if systems do not have downtime.
 - How do you do “immediate switchover”?
 - Hard to recover from failure, may require reverse data migration.

Question 6

Describe advantages and disadvantages for each technique in terms of data migration, complexity, and rollback in case of failure.

- **Parallel Run:**
 - Data migration happens in parallel with older system running (if new system fails on certain queries, route to old system).
 - Less risk in data migration. Need to have support for continuous migration as data is added to new system.
 - More complex than big bang.
 - Policies needed for synchronizing and routing traffic between old/new system.
 - How is the new system validated?
 - Need to maintain redundant copies of data.
 - Rollback can be run in lockstep.
 - Can also run split-stream traffic, then reverse data migration may be required.

Question 6

Describe advantages and disadvantages for each technique in terms of data migration, complexity, and rollback in case of failure.

- Staged:
 - Data migration needs facades to allow new system components to work with old system components.
 - These facades translate data back and forth.
 - Since you need facades, there is some complexity.
 - If migrating parts of organization, you may also need policies for handling discrepancies between the new/old system.

Question 6

Describe installation scenarios where you might use each technique.

- **Big bang:** Home web server upgrade
 - Just move everything over and be done.
 - Only affects you.
- **Parallel run:** Banking transaction software
 - Can test new system with a subset of users.
 - If any issues, use the existing reliable system.
- **Staged migration:** Air traffic control
 - Hard to keep old and new synchronized, but want to control release of new system to ensure each component is reliable.

Question 7

What are the availability benefits and risks associated with the following architectural styles: Pipe and Filter, Repository, Event-based, Layered.

Question 7

What are the availability benefits and risks associated with the following architectural styles: Pipe and Filter.

- Can duplicate streams of data to multiple CPUs transparently (from filter perspective).
 - If failure occurs, can redirect to a good stream.
- Can easily introduce “voters” to look for disagreements in results from multiple filters.
- Risks: On a single pipe & filter chain, any single failure will cause whole system to fail because filters do not know about each other.

Question 7

What are the availability benefits and risks associated with the following architectural styles: Repository.

- Risky for availability.
 - Central store, so consistency problems in presence of multiple readers/writers.
 - Central point of failure.
- Consolidates critical data in a single location, so single-point logging & recovery.
- Many schemes for high-availability repositories (database clusters).
 - Expensive, but work well in practice.

Question 7

What are the availability benefits and risks associated with the following architectural styles: Event-based.

- If there is a centralized event broker, can be risky for availability (and performance) because it introduces a central point of failure.
- Also, can be difficult to understand the composite behavior of event-based systems.
 - “Event storms” can occur if one event leads to a cascade of many events that can reduce system reliability and availability.
- On the other hand, failover mechanisms when constructing highly-available systems are event-based.
 - A heartbeat event, sent at regular intervals, is the means by which system health is monitored. If a sibling system does not send a heartbeat, then failover is performed.

Question 7

What are the availability benefits and risks associated with the following architectural styles: Layered.

- Layers help create highly available systems because they limit the kinds of failures that must be accounted for.
 - Web is reliable because of isolated handling of classes of failure by different layers. IP handles routing, TCP handles packet retransmission, load balancing handles server failures, etc.

Question 8

Write the guarantees for the following AADL component describing a dishwasher mode controller:

```
system Dishwasher_Mode
  features
    door_closed: in data port Base_Types::Boolean;
    time_remaining: in data port Base_Types::Integer;
    pump_on: out data port Base_Types::Boolean;
    dishwasher_mode: out data port Base_Types::Integer;

  annex agree {**
    const SETUP_MODE : int = 0;
    const WASHING_MODE : int = 1;
    const RINSE_MODE : int = 2;
    const DRYING_MODE : int = 3;

    guarantee "the pump shall be off if the door is open" : true;
    guarantee "If the dishwasher was in WASHING_MODE and time
remaining is zero, it shall enter RINSE_MODE" : true;
    guarantee "The dishwasher shall never transition directly from
WASHING_MODE to DRYING_MODE" : true;
    guarantee "The dishwasher shall start in SETUP_MODE" : true;
  **};
end Dishwasher_Mode;
```

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    guarantee "the pump shall be off if the door is open" :
      pump_on => door_closed;
      (alternate): (not door_closed) => (not pump_on);

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    guarantee "If the dishwasher was in WASHING_MODE and
time remaining is zero, it shall enter RINSE_MODE" :
      true -> ((pre(dishwasher_mode) = WASHING_MODE and
time_remaining = 0) => dishwasher_mode = RINSE_MODE)
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    directly from WASHING_MODE to DRYING_MODE" :
      true -> (pre(dishwasher_mode) = WASHING_MODE => (not
dishwasher_mode = DRYING_MODE))
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    const DRYING_MODE : int = 3;

    guarantee "The dishwasher shall start in SETUP_MODE" :
      dishwasher_mode = SETUP_MODE -> true.

  **};
end Dishwasher_Mode;
```

Question 9

- **What is the difference between response time and throughput?**
- **Give an example of a system with excellent throughput but poor response time and vice versa.**

Question 9

What is the difference between response time and throughput?

- Response time is from the client's perspective.
 - How long does it take to service my request?
- Throughput is from the server's perspective.
 - How many requests can be processed in a given time period?

Question 9

Give an example of a system with excellent throughput but poor response time and vice versa.

- A pipelined system may involve several processors working in tandem to solve a particular problem.
 - It may be able to process very large volumes of transactions (high throughput) due to partitioning the problem into segments that are handled sequentially, while still exhibiting poor response time.
 - (each segment takes time t , with number of segments s , so the total response time is $n*s$).
- Instead, imagine a single processor non-pipelined system that processes requests sequentially.
 - If there are few requests, it will have better response time than the pipelined system because there is no latency in servicing the request.
 - However, it will have very poor throughput under heavy load.

Question 10

What is the distinguishing characteristic of a *real time*, as opposed to a *non-real time* system? What is the difference between hard and soft real time systems?

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What is the distinguishing characteristic of a *real time*, as opposed to a *non-real time* system? What is the difference between hard and soft real time systems?

- For a real-time system, an operation's correctness depends not only on logical correctness, but the time required to complete it.
- In a hard real-time system, computation of an answer after its deadline is considered failure.
- Soft real time systems can tolerate missed deadlines as long as there is a bound on the number of missed deadlines within some time scale.

**Thanks for a great
semester!!!!!!!!!!!!!!!!!!!!!!**

Next Time

- The Final
 - Please e-mail questions + office hours today.
 - Exam will be proctored by a student, so ask me questions before the day of the test.

- Homework:
 - Project, Part 4 - Due tonight
 - Assignment 3 - Due on December 9