





Gregory Gay DIT635 - February 19, 2020



Test Adequacy Criteria

Instead - can we compromise between the impossible and the inadequate?



- Can we measure "good testing"?
- Test adequacy criteria "score" testing efforts by measuring the completion of test obligations.
 - Checklists of properties that must be met by test cases.





Structural Coverage Criteria

- Criteria based on exercising:
 - Statements (nodes of CFG)
 - Branches (edges of CFG)
 - Decisions and Conditions
 - Paths
 - ... and many more
- Measurements used as adequacy criteria

Elements Vs. Paths

- Statement, Branch, Condition Coverage all focus on one element at a time.
- A test executes a path, not a single element.
- Each element on that path is dependent on the others.

```
boolean A = ...
boolean B = ...
boolean expr = A || B;

if (expr && C) {
    System.out.println("Here I am!");
}
```





Elements Vs. Paths

- There are different control paths through a program...
- ... And different ways that data passed along paths can influence execution.
- Important to examine not just elements, but paths.

```
boolean A = ... Fault in definition
boolean expr = A || B; Corrupts definition
of expr if B = False

if (expr && C) {
    System.out.println("Here I am!");
}
expr can corrupt outcome if C = True
```

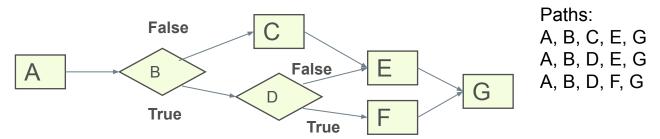


Today's Goals

- Introduce Path Coverage
- Data Flow Coverage Criteria
 - Focus on how information spreads through a program instead of how control flows.
 - Based on Definition-Use Pairs
 - (Where is X defined? Where is each definition of X used?)



 Path coverage requires that all paths through the CFG are covered.



Coverage = Number of Paths Covered

Number of Total Paths



Path coverage is a powerful coverage metric, but is often impractical.

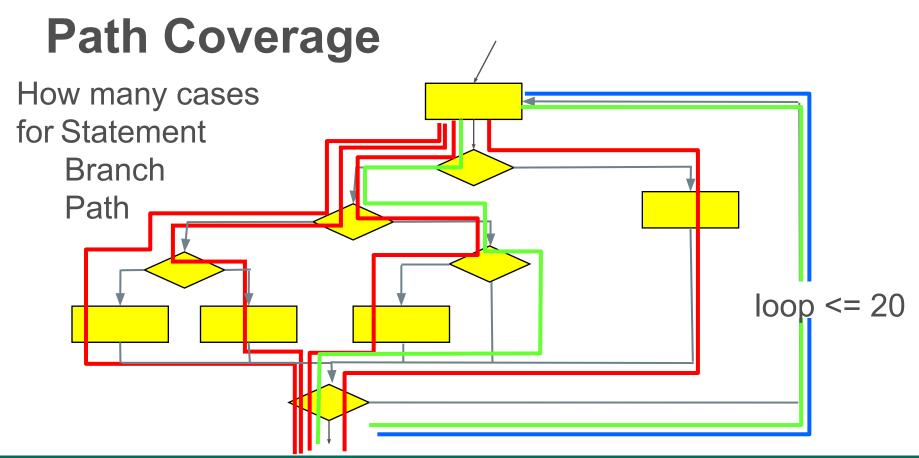
- How many paths does this have?
- Each loop cycle is a separate path!

True

A[i] = -A[i];

True





Path coverage with (loop <= 20) requires: **3,656,158,440,062,976** test cases

If you run 1000 tests per second, this will take 116,000 years.

However, there are ways to get some of the benefits of path coverage without the cost...



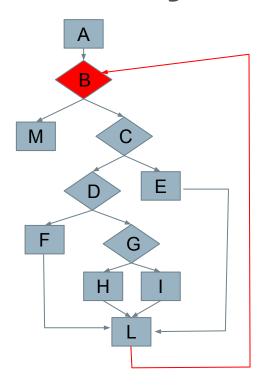
- Theoretically, a very strong coverage metric.
 - Many faults emerge through sequences of interactions.
- But... Generally impossible to achieve.
 - Loops result in an infinite number of path variations.
 - Even bounding number of loop executions leaves an infeasible number of tests.

Boundary Interior Coverage

- Groups paths that differ only in the subpath they follow when repeating the body of a loop.
 - Executing loop 20 times is different than executing it twice, but same *subpaths* repeat over and over.
 - Unroll loop in CFG into distinct subpaths, and cover those instead of worrying about loop cycles.



Boundary Interior Coverage



$$A -> B -> C -> D -> F -> L -> B$$



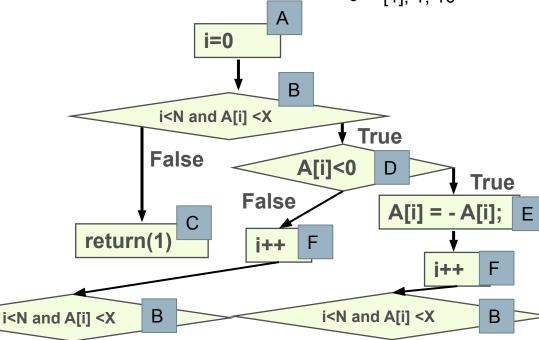
Boundary Interior Coverage

Paths:

- A, B, C
- A, B, D, F, B
- A, B, D, E, F, B

Tests

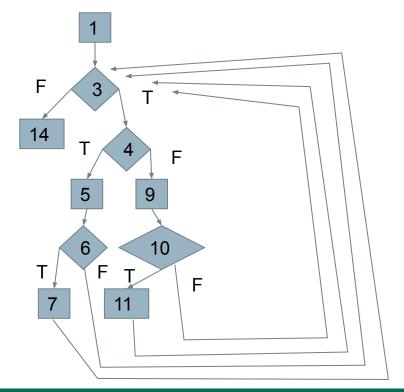
- [], 0, 10
- [-1], 1, 10
- [1], 1, 10





Boundary Interior Example

```
1. public int doSomething(int x, int y)
2. {
3.
     while(y > 0) {
4.
           if(x > 0) {
5.
                 y = y - x;
6.
                 if (y > 0)
                     System.out.println("Y: " + y);
7.
8.
          }else {
9.
                 x = x + 1;
10.
                 if (x \leftarrow 0)
11.
                     System.out.println(X: " + x);
12.
13.
14.
     return x + y;
15. }
```





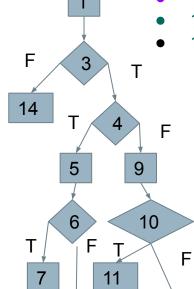
Boundary Interior Example

```
    public int doSomething(int x, int y)

2. {
3.
     while(y > 0) {
4.
           if(x > 0) {
5.
                y = y - x;
6.
                if (y > 0)
                     System.out.println("Y: " + y);
7.
8.
          }else {
9.
                x = x + 1;
10.
                if (x <= 0)
11.
                     System.out.println(X: " + x);
12.
13.
14.
     return x + y;
15. }
```

Paths:

- 1, 3-F, 14
- 1, 3-T, 4-T, 5, 6-T, 7, 3
- 1, 3-T, 4-T, 5, 6-F, 3
- 1, 3-T, 4-F, 9, 10-T, 11, 3
- 1, 3-T,4-F, 9, 10-F, 3



Tests:

- 10, -1
- 3, 4
- -1, 1

Number of Paths

- Boundary Interior Coverage removes bounds number of loop paths.
 - However, number of paths can still be exponential.
 - N non-loop branches results in 2^N paths.
- Additional limitations may need to be imposed.

```
if (a) S1;
if (b) S2;
if (c) S3;
...
if (x) SN;
```

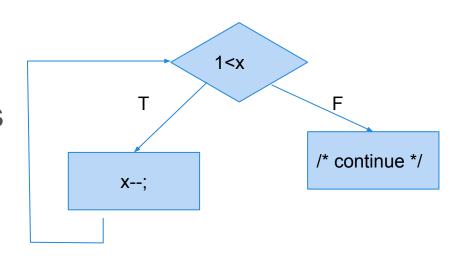
Data Flow





Control Flow

- Capture how execution navigates between blocks of statements.
- We care about a statement's effect only when it affects the path.
 - Deemphasizes information being transmitted.





Data Flow

- Program statements compute and transform data...
- Reason about data dependence
 - A variable is used here.
 - Where does its value come from?
 - Is this assigned value ever used?
 - Is this variable properly initialized?
 - If the expression assigned to a variable is changed what else would be affected?



Data Flow

- Basis of the optimization performed by compilers.
- Used to derive test cases.
 - Have we covered the dependencies?
- Used to detect faults and other anomalies.
 - When can we cache result of a calculation instead of recalculating it?
 - Can we eliminate a variable definition?



Definition-Use Pairs

- Data is defined.
 - ... and data is used.
- Pairs of definitions and uses capture flow of information through the program.
 - Definitions occur when variables are declared, initialized, assigned values, or received as parameters.
 - Uses occur in expressions, conditional statements, parameter passing, return statements.

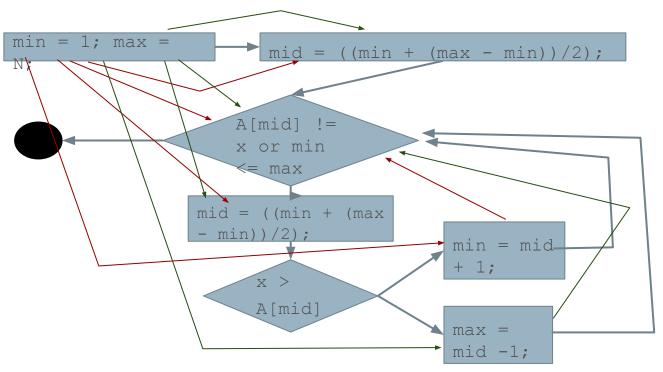
Definitions and Uses

```
1. min = 1;
    max = N;
    mid = ((min + (max - min))/2);
    while (A[mid] != x or min <= max){
5.
        mid = ((min + (max - min))/2);
6.
         if (x > A[mid]){
             min = mid + 1
7.
8.
         } else {
             max = mid - 1;
9.
10.
11.
```

- 1. **def** min
- 2. def max, use N
- 3. **def** mid, **use** min, max
- 4. **use** A[mid], mid, x, min, max
- 5. **def** mid, **use** min, max
- 6. **use -** x, A[mid], mid
- 7. **def -** min, **use** mid
- 8. -
- 9. def max, use mid



Definitions and Uses



- 1. **def min**
- 2. def max, use N
- 3. **def** mid, **use** min, max
- 4. **use** A[mid], mid, x, min, max
- 5. **def** mid, **use** min, max
- **6. use** x, A[mid], mid
- 7. def min, use mid
- 8. -
- 9. def max, use mid



Definition-Use Pairs

- We can say there is a def-use pair when:
 - There is a **def** (definition) of variable x at location A.
 - Variable x is used at location B.
 - A control-flow path exists from A to B.
 - and the path is **definition-clear** for x from A to B.
- If a variable is redefined, the original def is **killed** and the pair is between the new definition and its use in B.

Example - Definition-Use Pairs

```
1. min = 1;
2. max = N;
    mid = ((min + (max - min))/2);
    while (A[mid] != x or min <= max){
        mid = ((min + (max - min))/2);
 5.
6.
        if (x > A[mid]){
7.
            min = mid + 1
8.
        } else {
9.
            max = mid - 1;
10.
11.
```

DU Pairs min: (1, 3), (1, 4), (1, 5), (7, 4), (7, 5)max: (2, 3), (2, 4), (1, 5), (9, 4), (9, 5)N: (0, 2) mid: (3, 4), (5, 6), (5, 7), (5, 9), (5, 4)x: (0, 4), (0, 6) A: (0, 4), (0, 6)

Example - GCD

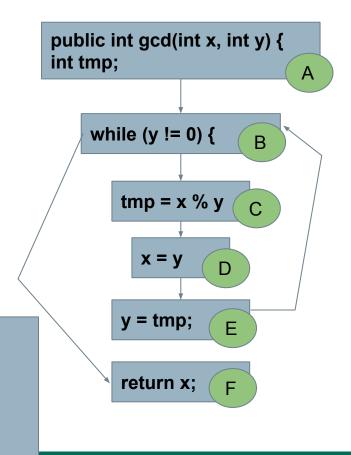
```
public int gcd(int x, int y){
2.
        int tmp;
3.
        while(y!=0){
4.
            tmp = x \% y;
5.
           X = y;
6.
           y = tmp;
7.
8.
        return x;
```

- 1. def: x, y
- 2. def: tmp
- 3. use: y
- 4. use: x, y def: tmp
- 5. use: y def: x
- 6. use: tmp def: y
- 7. -
- 8. use: x



Example - GCD

```
public int gcd(int x, int y){
2.
         int tmp;
3.
         while(y!=0){
              tmp = x \% y;
4.
5.
              X = y;
6.
              y = tmp;
7.
8.
         return x;
9.
        Def-Use Pairs
        x: (1, 4), (5, 4), (5, 8), (1, 8)
        y: (1, 3), (1, 4), (1, 5), (6, 3), (6, 4), (6, 5)
        tmp: (4, 6)
```





Example - collapseNewlines

```
7. public static String collapseNewlines(String argStr)
8. {
9.
     char last = argStr.charAt(0);
10.
     StringBuffer argBuf = new StringBuffer();
11.
     for(int cldx = 0; cldx < argStr.length(); cldx++)</pre>
12.
13.
14.
           char ch = argStr.charAt(cldx);
15.
           if(ch != '\n' || last != '\n')
16.
17.
                 argBuf.append(ch);
18.
                 last = ch;
19.
20.
21.
22.
     return argBuf.toString();
23. }
```

Variable	D-U Pairs
argStr	(7, 9), (7,12), (7, 14)
last	(9, 15), (18, 15)
argBuf	(17, 22)
cldx	(12, 12), (12, 14)
ch	(14, 15), (14, 17), (14, 18)

Let's Take a Break

Dealing With Arrays and Pointers

- Arrays and pointers (including object references and arguments) introduce issues.
 - It is not possible to determine whether two access refer to the same storage location.

```
    a[x] = 13;
    k = a[y];
    Are these a def-use pair?
```

- a[2] = 42; i = b[2];
 - Are these a def-use pair?
 - Aliasing = two names refer to the same memory location.

Aliasing

- Aliasing is when two names refer to the same memory location.
 - int[] a = new int[3]; int[] b = a;
 a[2] = 42;
 i = b[2];
 - a and b are aliases.
- Worse in C:

$$p = &b$$
* $(p + i) = k;$



Uncertainty

- Aliasing introduces uncertainty.
 - Instead of definition or use of one variable, may have a potential def or use of a set of variables.
- Proper treatment depends on purpose of analysis:
 - Safest method treat any use of a potential alias of V as a use of V.
 - Creates more def-use pairs (some may not be real pairs), but avoids missed pairs.

Dealing With Uncertainty

Treat all potential aliases as definitions and uses:

```
a[1] = 13;
k = a[2];

Def of a[1], use of a[2].

a[x] = 13;
a[x] = a[y];

Def and use of array a.
```

- · Can be very imprecise.
 - They are only the same if x and y are the same.



Dealing With Uncertainty

 Option 2: Treat uncertainty about aliases like uncertainty about control flow.

```
a[x] = 13;

k = a[y];

a[x] = 13;

if(x == y)

k = a[x];

else

k = a[y];
```

- Rewrite code to make references explicit.
- In transformed code, all array references are distinct.



Situational Def-Use Pairs

- ++counter, counter++, counter+=1
 counter = counter + 1
 - Use of counter then a new definition.
- char *ptr = *otherPtr
 - Definition of string *ptr
 - Use of memory index ptr, string *otherPtr, and memory index otherPtr.
 - ptr++
 - Use of memory index ptr, definition of both memory index and string *ptr (change to index moves pointer to a new location).

Dealing With Nonlocal Information

- fromCust and toCust may be references to same object.
 - from/toHome and from/toWork.
- Option 1 treat all nonlocal variables of same type as potential aliases.
- Option 2 Introduce additional control flow
 - if (fromHome.equals(fromWork))

```
public void transfer(Customer fromCust,
Customer toCust){
    PhoneNum fromHome =
         fromCust.getHomePhone();
    PhoneNum fromWork =
         fromCust.getWorkPhone();
    PhoneNum toHome =
         toCust.getHomePhone();
    PhoneNum toWork =
         toCust.getWorkPhone();
}
```

Data Flow Coverage Criteria

Overcoming Limitations of Path Coverage

- We can potentially expose many faults by targeting particular paths of execution.
- What are the important paths to cover?
 - Some methods impose heuristic limitations.
 - Use data flow to select paths based on how one element can affect the computation of another.



Choosing the Paths

- Computing the wrong value leads to a failure only when that value is used.
 - Pair definitions with usages.
 - Ensure that definitions are actually used by covering paths from definitions to uses.
 - All DU Pair Coverage, All DU Paths Coverage, All Definitions Coverage
 - Varying power and cost.



- Requires each DU pair be exercised in at least one program execution.
 - Cover any path between a definition and its use.
- Coverage = number exercised DU pairs
 number of DU pairs
- Can easily achieve structural coverage without covering all DU pairs.



All DU Pairs Coverage Example

```
    public int doSomething(int x, int y)

2. {
3.
     while(y > 0) {
4.
           if(x > 0) {
5.
                 y = y - x;
6.
                 if (y > 0)
                     System.out.println("Y: " + y);
7.
8.
          }else {
9.
                 x = x + 1;
10.
                 if (x \leftarrow 0)
                     System.out.println(X: " + x);
11.
12.
13.
14.
    return x + y;
15. }
```

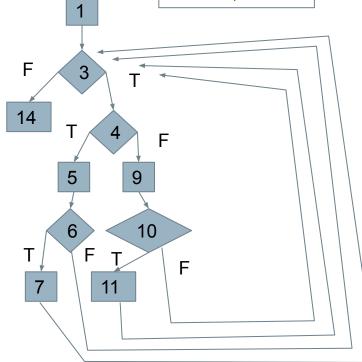
```
X:
(1, 4), (1, 5), (1, 9), (1, 14)
(9, 10), (9, 11), (9, 5), (9, 9), (9, 14)
Y:
(1, 3), (1, 5), (1, 14)
(5, 6), (5, 7), (5, 3), (5, 5), (5, 14)
```

```
X: (1, 4), (1, 5), (1, 9), (1, 14), (9, 10), (9, 11), (9, 5), (9, 9), (9, 14)
Y: (1, 3), (1, 5), (1, 14), (5, 6), (5, 7), (5, 3), (5, 5), (5, 14)
```

```
1. public int doSomething(int x, int y)
2. {
3.
     while(y > 0) {
4.
           if(x > 0) {
5.
                y = y - x;
6.
                if (y > 0)
                    System.out.println("Y: " + y);
7.
8.
          }else {
9.
                x = x + 1;
10.
                if (x <= 0)
                    System.out.println(X: " + x);
11.
12.
13.
14.
     return x + y;
15. }
```

Tests:

- 1. -1, 1
- 2. 3, 7
- 3. -2, 1





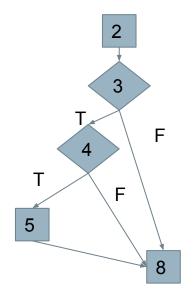
All DU Paths Coverage

- A use may be reachable along several paths from the definition.
- Cover all simple (non-looping) paths at least once.
 - Can reveal faults where a path is exercised that should use a certain definition but doesn't.

Coverage = number of exercised DU paths number of DU paths

All DU Paths Example

```
1. ...
2. int x = 1;
3. if (y > 7) {
4.    if (z > 5) {
5.       z = x + 5;
6.    }
7. }
8. y = x + 7;
9. ...
```



DU Pair (2, 8) for X can be reached along multiple paths.

- 2, 3T, 4T, 5, 8
- 2, 3T, 4F, 8
- 2, 3F, 8

Tests:

- y = 10, z = 6
- y = 10, z = 3
- y = 2, z = (anything)

Path Explosion Problem

- Even without looping paths, number of DU paths can be exponential.
 - Code between definition and use can be irrelevant to that variable, but contains many paths.

```
public void countBits(char ch){
    int count = 0;
    if (ch & 1)
                  ++count;
    if (ch & 2)
                  ++count;
    if (ch & 4)
                  ++count;
    if (ch & 8)
                  ++count;
    if (ch & 16)
                  ++count;
    if (ch & 32)
                  ++count;
    if (ch & 64)
                  ++count;
    if (ch & 128) ++count;
    System.out.println(ch + " has " +
count + "bits set to 1");
```



All Definitions Coverage

- All DU Pairs/All DU Paths may be too expensive in some situations.
- Pair each definition with at least one use.
 - Skips many DU pairs, but ensures each definition tried.

Coverage = number of covered definitions

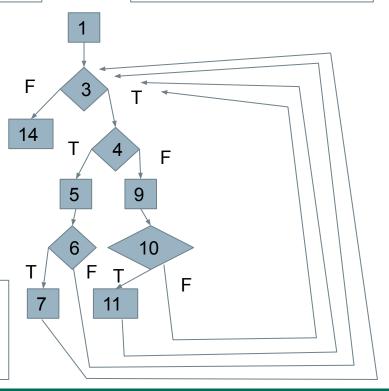
number of definitions

```
X: (1, 4), (1, 5), (1, 9), (1, 14), (9, 10), (9, 11), (9, 5), (9, 9), (9, 14)
Y: (1, 3), (1, 5), (1, 14), (5, 6), (5, 7), (5, 3), (5, 5), (5, 14)
```

X: Definitions on lines 1, 9 Y: Definitions on lines 1, 5

```
1. public int doSomething(int x, int y)
2. {
3.
     while(y > 0) {
4.
           if(x > 0) {
5.
                 y = y - x;
6.
                 if (y > 0)
                     System.out.println("Y: " + y);
7.
8.
           }else {
9.
                 x = x + 1;
10.
                 if (x \leftarrow 0)
11.
                     System.out.println(X: " + x);
12.
13.
14.
     return x + y;
15. }
```

- Any test covers (1, -) pairs.
- Reaching lines 5, 9 gets (5,14) and (9,14) pairs.





Infeasibility Problem

- Metrics may ask for impossible test cases.
- Path-based metrics may require infeasible combinations of feasible elements.
 - Alias analysis may add additional infeasible paths.
- All Definitions, All DU-Pairs Coverage reasonable.
 - All DU-Paths is much harder!



Activity - DU Pair Coverage

- Identify all DU pair
- Write test cases to achieve All DU Pair Coverage.
 - Hint remember that there is a loop.

https://bit.ly/2No5qqM

```
1. int doSomething(int x, int y)
2. {
        while(y > 0) {
3.
            if(x > 0) {
5.
                y = y - x;
6.
            }else {
7.
                x = x + 1;
8.
9.
10.
        return x + y;
11. }
```



Activity - DU Pairs

```
1. int doSomething(int x, int y)
2. {
3.
        while(y > 0) {
            if(x > 0) {
4.
5.
                y = y - x;
6.
            }else {
7.
                x = x + 1;
8.
9.
10.
        return x + y;
11. }
```

Variable	Defs	Uses
х	1, 7	4, 5, 7, 10
у	1, 5	3, 5, 10

Variable	D-U Pairs	
х	(1, 4), (1, 5), (1, 7), (1, 10), (7, 4), (7, 5), (7, 7), (7, 10)	
у	(1, 3), (1, 5), (1, 10), (5, 3), (5, 5), (5, 10)	



Activity - DU Pairs

```
1. int doSomething(int x, int y)
2. {
3.
        while(y > 0) {
4.
             if(x > 0) {
5.
                 y = y - x;
             }else {
6.
7.
                 x = x + 1;
8.
9.
10.
        return x + y;
11. }
```

Variable	Defs	Uses
x	1, 7	4, 5, 7, 10
у	1, 5	3, 5, 10

Variable	D-U Pairs
х	(1, 4), (1, 5), (1, 7), (1, 10), (7, 4), (7, 5), (7, 7), (7, 10)
у	(1, 3), (1, 5), (1, 10), (5, 3), (5, 5), (5, 10)

Test 1:
$$(x = 1, y = 2)$$

Covers lines 1, 3, 4, 5, 3, 4, 5, 3, 10
Test 2: $(x = -1, y = 1)$
Covers lines 1, 3, 4, 6, 7, 3, 4, 6, 7, 3, 4, 5, 3, 10
Test 3: $(x = 1, y = 0)$
Covers lines 1, 3, 8



We Have Learned

- Control-flow and data-flow capture important paths in program execution.
- Analysis of how variables are defined and then used and the dependencies between definitions and usages can help us reveal important faults.
- Many forms of analysis can be performed using data flow information.



We Have Learned

- If there is a fault in a computation, we can observe it by looking at where the computation is used.
- By identifying DU pairs and paths, we can create tests that trigger faults along those paths.
 - All DU Pairs coverage
 - All DU Paths coverage
 - All Definitions coverage



Next Time

- Exercise Session Structural Testing
 - Using Meeting Planner code.
- Next Wednesday Fault-Based Testing
 - Pezze & Young Ch 16
- Assignment 2
 - Due February 28! We have covered everything on it.



UNIVERSITY OF GOTHENBURG

